

THE FINAL SECTIONS

7. More on soundness

- more constraints on the letter-refiner (than just e-soundness):
conditional frame
soundness

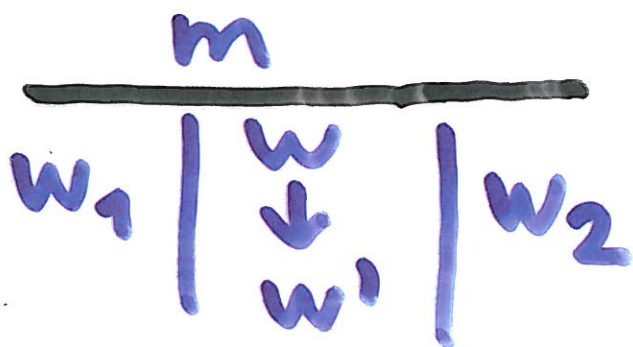
8. [Very briefly] Directions towards practical applications

- coalescence of the refinement contexts
- introducing coupled leaf expansion
- what about back-tracking?

9. summary

REWRITING MAPS

- a map is a set of 4-tuples of words
- $\langle w_1, w, w_2, w' \rangle \in m$ can be depicted as



- m is ϵ -sound if

$$w_1 w w_2 \stackrel{\epsilon}{\sim} w_1 w' w_2$$

for all 4-tuples

$$\langle w_1, w, w_2, w' \rangle \in m$$

LETTER-REFINER m -LEGITIMACY

- a letter-refiner r is m -legitimate if

$$w' \in r(w_1, c, w_2)$$

\Rightarrow

$$\frac{m}{w_1 \mid c \mid w_2} \\ w_1 \mid \downarrow \mid w_2 \\ w_1 \mid w' \mid w_2$$

- Note: we assume $r(w_1, a, w_2) = \{a\}$ for every terminal a

CONDITIONAL FRAME SOUNDNESS

- a frame is m-sound if every tetrasystem with the frame is e-sound whenever
 - ① the letter-refiner is m-legitimate
 - and
 - ② m is e-sound
- Let M be a set of maps. We say a frame is M-sound if it is m-sound for every $m \in M$
- We are interested in classes consisting of all maps with particular closure properties ...

MAP ADJUNCTIVITY

$$w_1 \left| \begin{array}{c} w^1 \\ \downarrow \\ w'' \end{array} \right| w^x w_2 \wedge w_1 w^1 \left| \begin{array}{c} w^x \\ \downarrow \\ w^{xx} \end{array} \right| w_2$$

$$\Rightarrow w_1 \left| \begin{array}{c} w^1 w^x \\ \downarrow \\ w'' w^{xx} \end{array} \right| w_2$$

- Let M_A denote the class of all adjunctive maps
- for example:

$$\Delta \left| \begin{array}{c} A \\ \downarrow \\ P \end{array} \right| B \wedge A \left| \begin{array}{c} B \\ \downarrow \\ Q \end{array} \right| \Delta$$

$$\Rightarrow \Delta \left| \begin{array}{c} AB \\ \downarrow \\ PQ \end{array} \right| \Delta$$

Theorem:

$$\begin{aligned} & \langle \theta_c, \theta_c, \theta_c, \theta_c \rangle \\ & \langle \theta_c, \theta_c, \Delta \theta_c, \Delta \theta_c \rangle \\ & \langle \theta_c, \theta_c, \theta_E, \theta_E \rangle \end{aligned}$$

are all M_A -sound

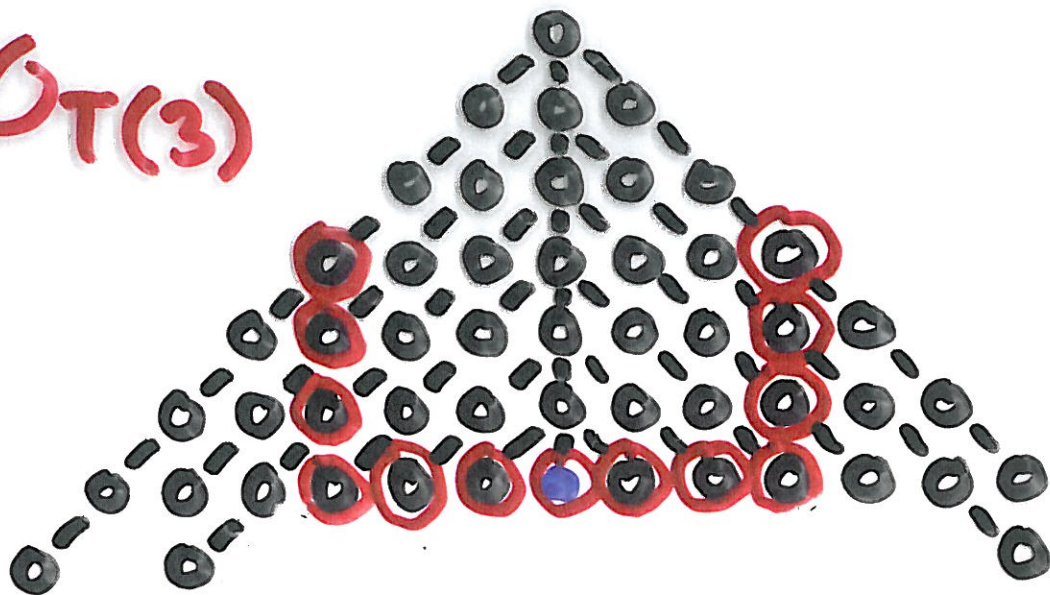
MAP TRANSITIVITY

$$w_1 \begin{array}{c} | \\ \downarrow \\ \end{array} \begin{array}{c} w \\ w' \\ \end{array} | w_2 \wedge w_1 \begin{array}{c} | \\ \downarrow \\ \end{array} \begin{array}{c} w' \\ w'' \\ \end{array} | w_2 \\ \Rightarrow w_1 \begin{array}{c} | \\ \downarrow \\ \end{array} \begin{array}{c} w \\ w'' \\ \end{array} | w_2$$

- Let M_{TA} denote the class of all such maps that are both transitive and adjunctive

Remark: This is where we cut a longer story short in this introductory presentation: we do not consider any other closure properties than these two

$\mathcal{G}_T(3)$



Let $1 \leq k < \infty$.

$$\varphi_{\mathcal{G}_T(k)}(i) = \min \{ |i|, k \}$$

Consider frame $(\mathcal{G}_T(k), \mathcal{G}_T(k), \mathcal{G}_E, \mathcal{G}_E)$

- it is confluent
- it is k -distributively progressive

Theorem:

- it is also M_{TA} -sound