

# TODAY's MENU

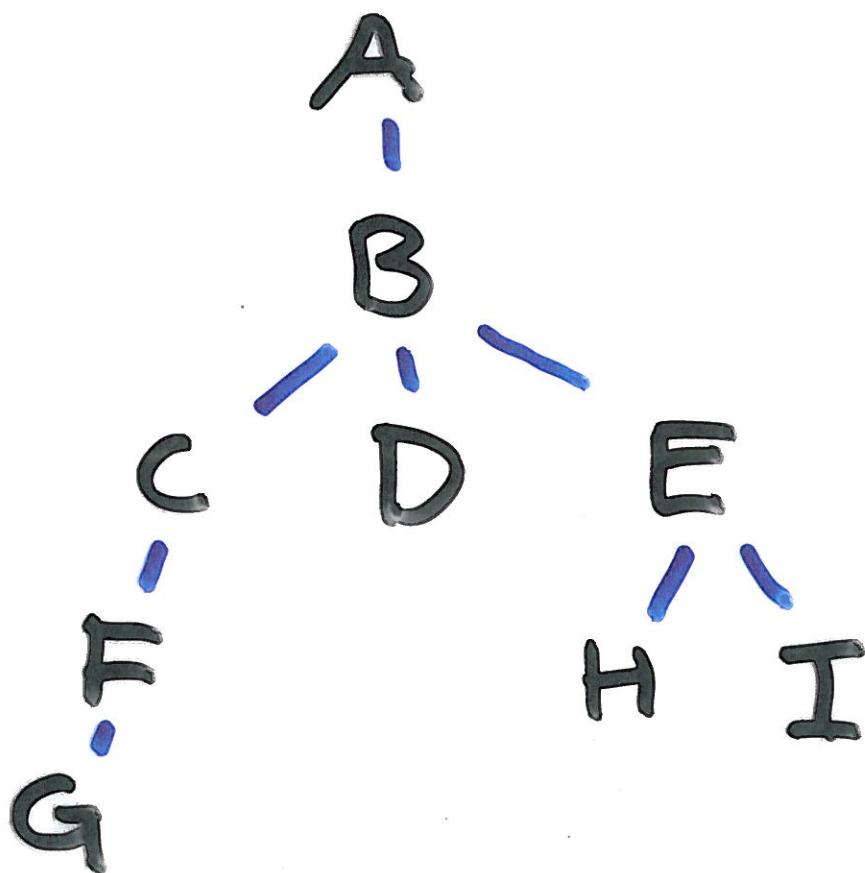
## 1. Introduction

## → 2. Belt-selectors

- belts
- angles
- (belt-assigners  
and) belt-  
selectors
- bijection  
between  
belt-selectors  
and combs
- examples

## 3. Tetrasytems

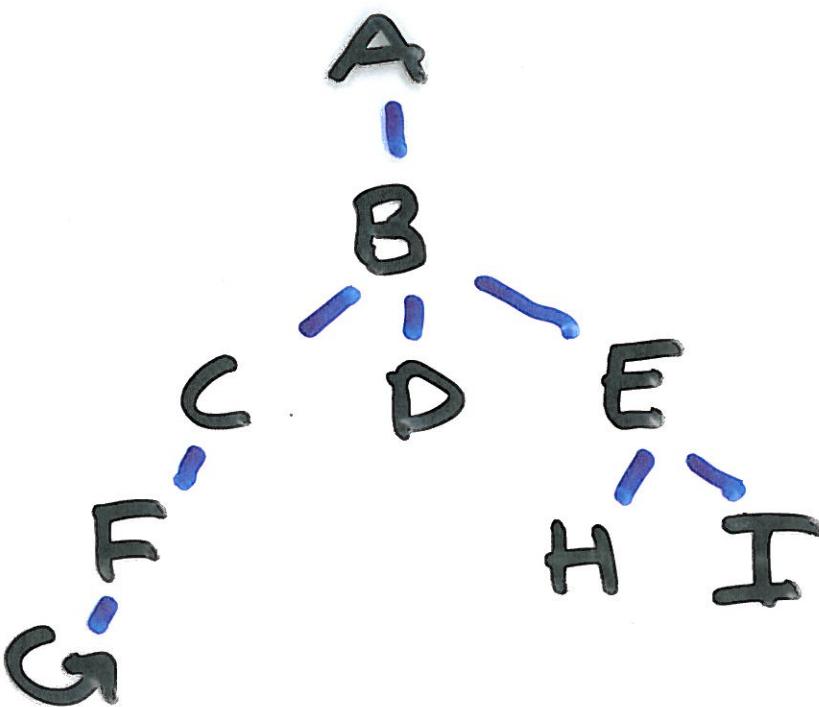
# BELTS



{A}, {B}, {C, D, E},  
{F, D, H, I}, {G, D, H, I},  
etc.

Each leaf has exactly  
one ancestor  
in the belt !

# ANGLES



$$\#(G, E) = \langle 3, 1, 1 \rangle$$
$$\#(E, G) = \langle 1, -1, 3 \rangle$$

$$\#(G, D) = \langle 3, 1, 1 \rangle$$
$$\#(D, G) = \langle 1, -1, 3 \rangle$$

$$\#(H, B) = \langle 2, 0, 0 \rangle$$
$$\#(B, H) = \langle 0, 0, 2 \rangle$$

An auxiliary definition:

a belt-provider is any function taking a tree and one of its nodes and returning a belt of the tree

- Like:

$$s(x, n) = \{n_1, n_2, n_3\}$$

# MORE INTERESTING DEFINITIONS

- a further condition on a belt-assigner

$$n^* \in s(x, n) \Leftrightarrow m^* \in s(y, m)$$

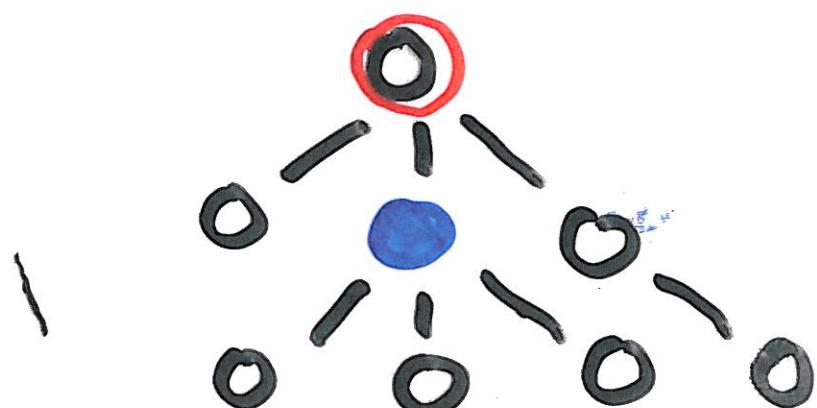
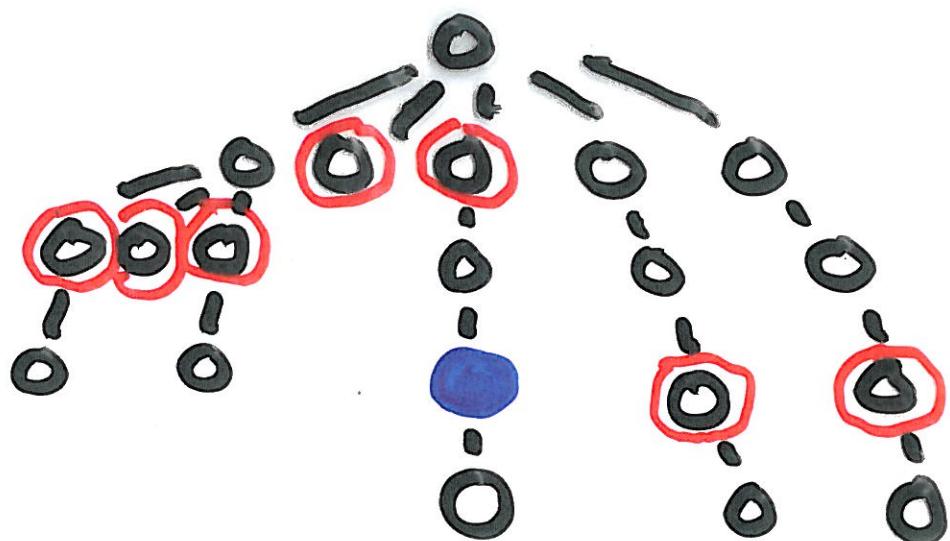
if all the following conditions are met:

- ①  $\not\exists (n, n^*) = \not\exists (m, m^*)$
- ②  $n^*$  is root  $\Leftrightarrow m^*$  is root
- ③  $n^*$  is leaf  $\Leftrightarrow m^*$  is leaf

- a belt-selector:

the above definition is tightened by dropping condition ②

# A SAMPLE BELT-ASSIGNMENT



HINT:

$$f(-3) = 2$$

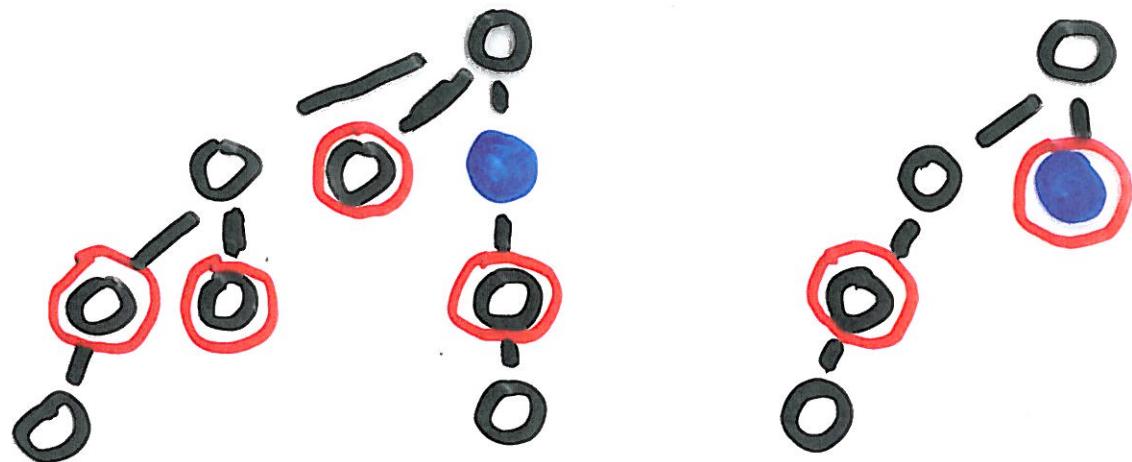
$$f(-2) = f(2) = 0$$

$$f(-1) = f(1) = -1$$

$$f(0) = -2$$

$$f(3) = 3$$

# A SAMPLE BELT-SELECTOR



$$f(-1) = 2$$
$$f(0) = 1$$

Note: a belt-selector never selects a proper ancestor of the argument node

[Now we forget belt-providers and belt-assigners]

# COMBS

- There is a bijection between belt-selectors and combs

A comb is any function

$$f: \{\dots, -2, -1, 0, 1, 2, \dots\} \rightarrow \{0, 1, 2, \dots\} \cup \{\infty\}$$

with  $f(i) = 0 \Rightarrow i = 0$

A given belt-selector  $s$ , whose comb is denoted as  $f$ , selects the following nodes for a given node  $n$

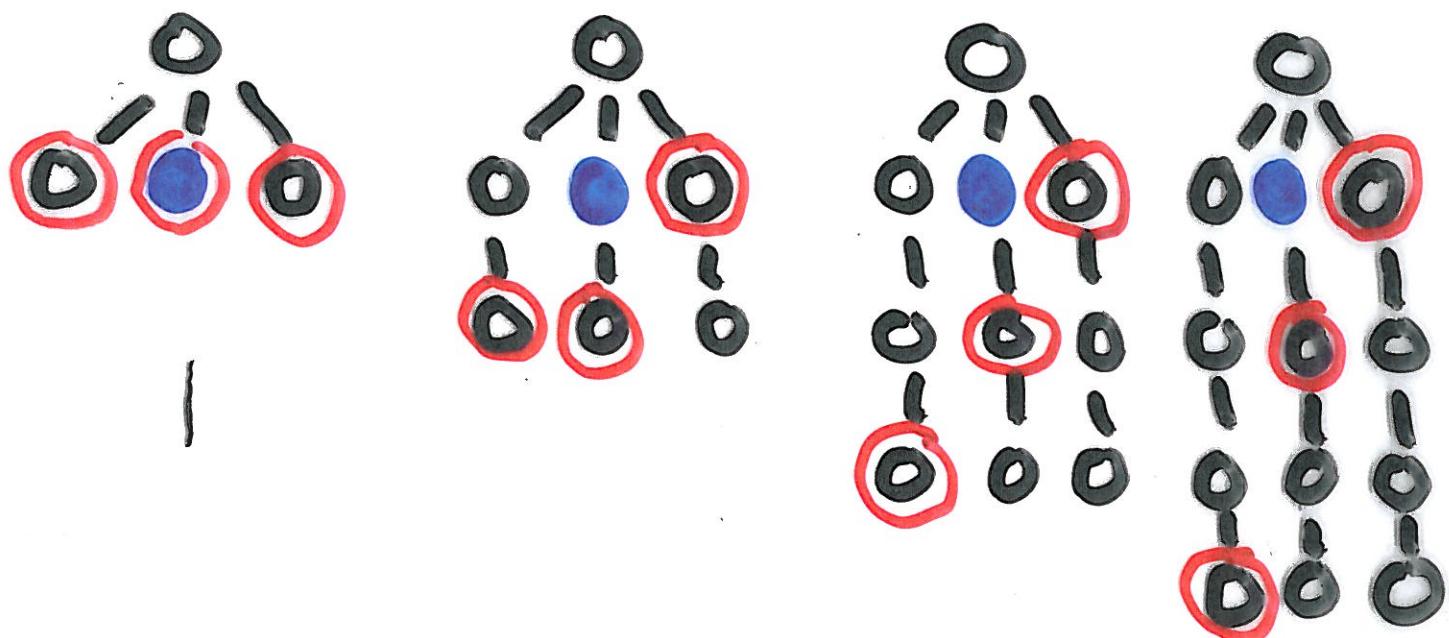
- ① all nodes  $n'$  with  
 $\chi(n, n') = \langle i, d, f(i * d) \rangle$
- ② all leaves with no ancestor already selected by ①

but not proper ancestors of  $n$ !!

# EXAMPLE

Suppose our belt-selector has the following comb :

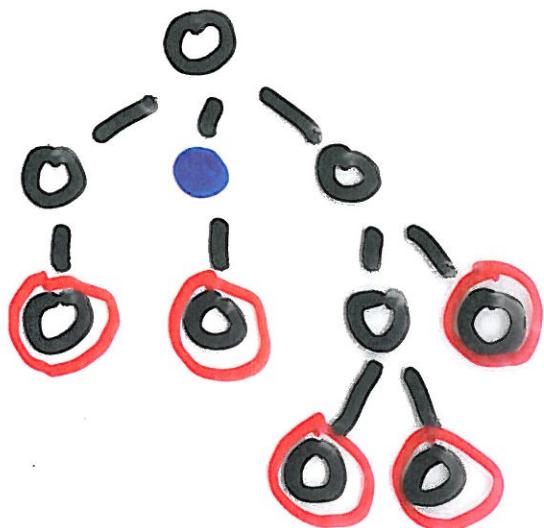
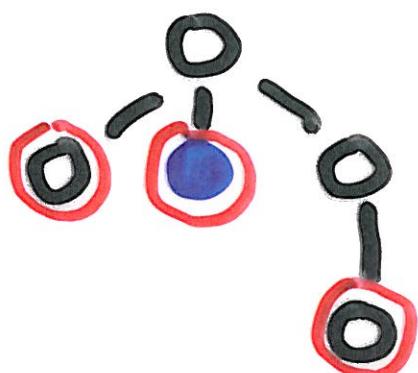
$$f(i) = \begin{cases} \infty & \text{when } i < 0 \\ 1 & \text{when } i \geq 0 \end{cases}$$



# SPECIFIC BELT-SELECTORS

## $\sigma_E$

$$\varphi_{\sigma_E}(i) = \infty$$

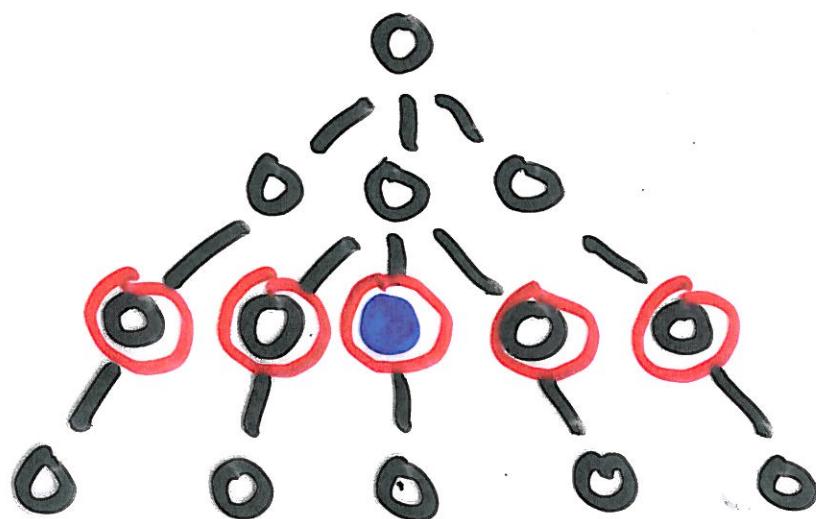


So  $\sigma_E$  always  
selects all  
the leaves!

# ANOTHER SPECIFIC BELT-SELECTOR:

$\sigma_C$

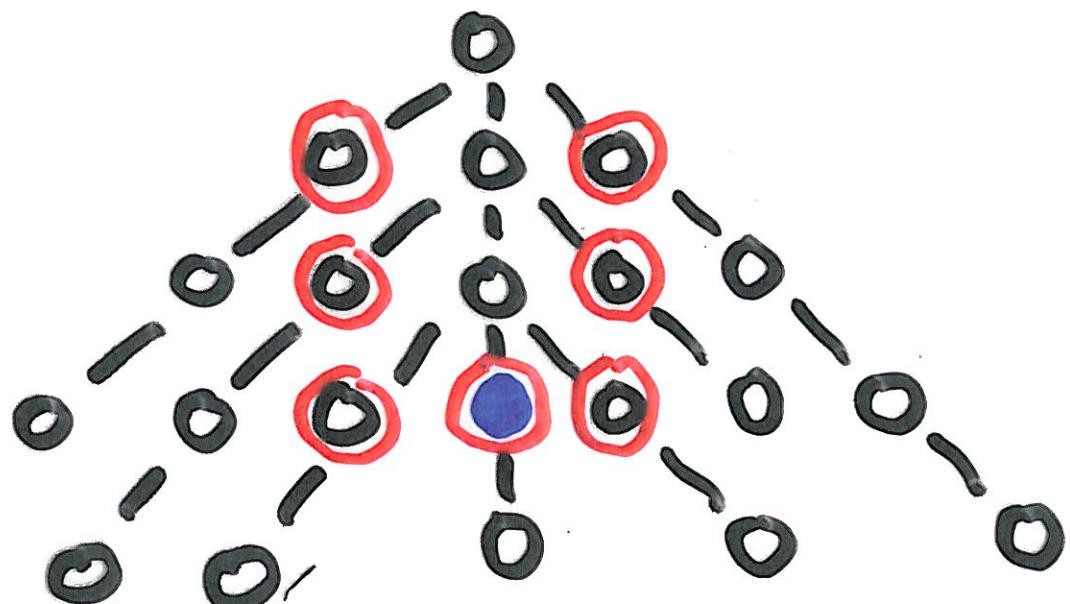
$$\varphi_{\sigma_C}(i) = |i|$$



# YET ANOTHER SPECIFIC BELT-SELECTOR:

$G_I$

$$\varphi_{G_I}(i) = \begin{cases} 0 & \text{when } i=0 \\ 1 & \text{when } i \neq 0 \end{cases}$$

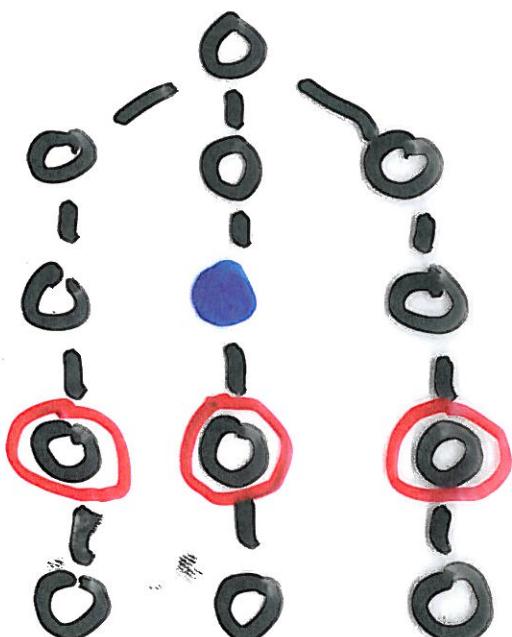
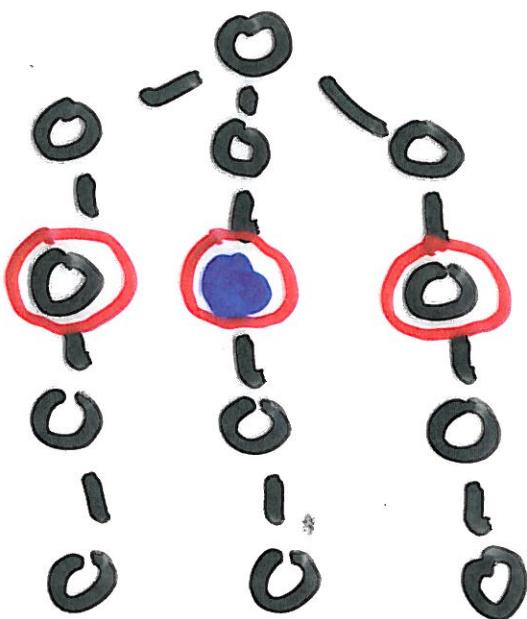


So  $G_I$  is the least greedy possible (while  $G_E$  is the most greedy)!

# ENTER AN OPERATOR :

$\Delta S$

$$\begin{aligned}\varphi_{\Delta G_C}(i) &= \varphi_{G_C}(i) + 1 \\ &= |i| + 1\end{aligned}$$



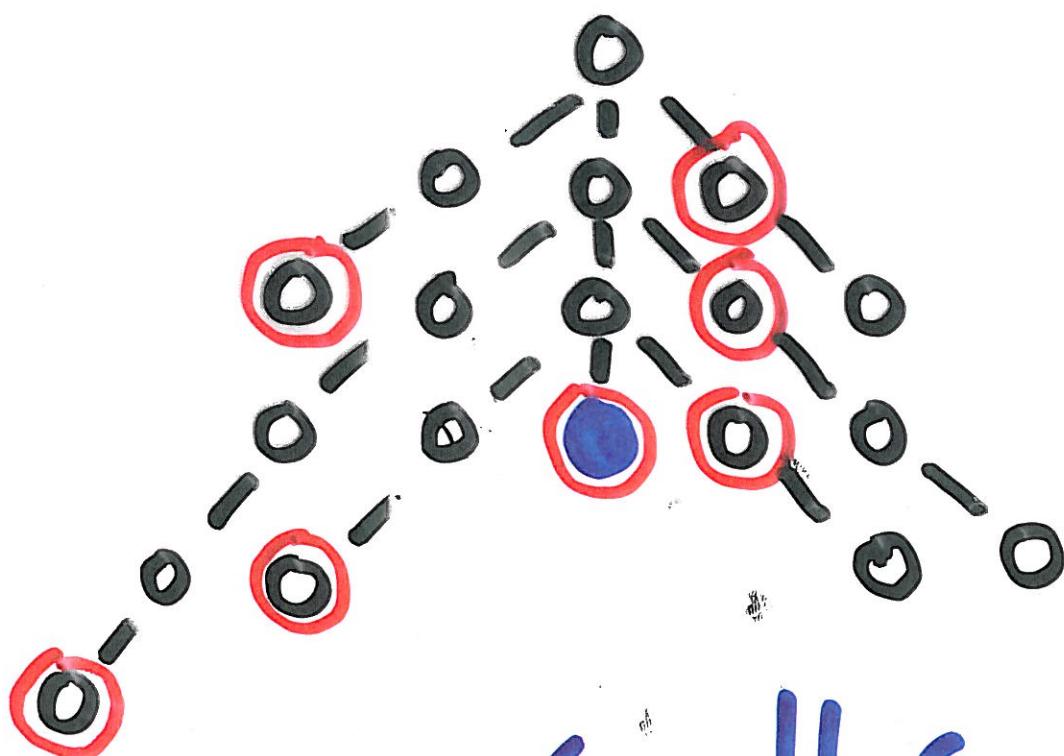
$G_C$

$\Delta G_C$

# FINALLY, A BINARY OPERATOR:

$s_1 \parallel s_2$

$$\varphi_{\delta_E \parallel \delta_I} = \begin{cases} \varphi_{G_E}(i) & \text{when } i < 0 \\ 0 & \text{when } i = 0 \\ \varphi_{G_I}(i) & \text{when } i > 0 \end{cases}$$



$\delta_E \parallel \delta_I$